

Some optimal codes related to graphs invariant under the alternating group A_8 .

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Primitive Rank-3 groups on Symmetric Designs

- In a classification paper [Dempwolff \(2001\)](#) determined the symmetric designs that admit a group which has a non-abelian socle and is primitive rank-3 on points and blocks.
- As a by product, the existence and uniqueness of a symmetric $2-(35, 17, 8)$ design having the simple alternating group A_8 as a non-abelian socle and acting primitively as rank-3 on points and blocks of the design was proved.
- This talk is about the structures related to to this design.

Preliminaries

- A result of Key and J Moori on designs, graphs and codes from primitive representation of a finite group outlines a construction of symmetric 1–designs

Result (1)

Let G be a *finite primitive permutation group* acting on the set Ω of size n . Let $\alpha \in \Omega$, and let $\Delta \neq \{\alpha\}$ be an orbit of the stabilizer G_α of α . If $\mathcal{B} = \{\Delta^g \mid g \in G\}$ and, given $\delta \in \Delta$, $\mathcal{E} = \{\{\alpha, \delta\}^g \mid g \in G\}$, then $\mathcal{D} = (\Omega, \mathcal{B})$ forms a *symmetric 1- $(n, |\Delta|, |\Delta|)$ design*. Further, if Δ is a *self-paired orbit* of G_α then $\Gamma = (\Omega, \mathcal{E})$ is a *regular connected graph* of valency $|\Delta|$, \mathcal{D} is self-dual, and G acts as an automorphism group on each of these structures, primitive on vertices of the graph, and on points and blocks of the design.

$t - (v, k, \lambda)$ Designs

- An incidence structure $\mathcal{D} = (\mathcal{P}, \mathcal{B}, \mathcal{I})$ with **point set** \mathcal{P} and **block set** \mathcal{B} and incidence $\mathcal{I} \subseteq \mathcal{P} \times \mathcal{B}$ is a $t - (v, k, \lambda)$ design if
 - ▶ $|\mathcal{P}| = v$;
 - ▶ every block $B \in \mathcal{B}$ is incident with precisely k points;
 - ▶ every t distinct points are together incident with precisely λ blocks. t, v, k and λ are non-negative integers;
 $|\mathcal{B}| = b$ is the number of blocks;
 r = replication number = number of blocks per point;
for $t = 2$, the order of \mathcal{D} is $n = r - \lambda$.

An **incidence matrix** for \mathcal{D} is a $b \times v$ matrix $A = (a_{ij})$ of 0's and 1's such that

$$a_{ij} = \begin{cases} 1 & \text{if } (p_j, B_i) \in \mathcal{I} \\ 0 & \text{if } (p_j, B_i) \notin \mathcal{I} . \end{cases}$$

The group A_8

- We consider G to be the simple alternating group A_8 .
- Notice that G is also the group of invertible 4×4 matrices whose determinant is 1, over \mathbb{F}_2 .

No.	Max. sub.	Degree	#	length		
1	A_7	8	2	7		
2	$2^3 : L_3(2)$	15	2	14		
3	$2^3 : L_3(2)$	15	2	14		
4	S_6	28	3	12	15	
5	$2^4 : (S_3 \times S_3)$	35	3	16	18	
6	$(A_5 \times 3) : 2$	56	4	10	15	30

Table: Orbits of the point-stabilizer of A_8

Graphs, Designs and Codes from the reprn of degree 35

- Observe from Table 1 that there is just one class of maximal subgroups of A_8 of index 35.
- The stabilizer of a point is a maximal subgroup isomorphic to the group $2^4 : (S_3 \times S_3)$. rank-3 primitive group on the cosets of $2^4 : (S_3 \times S_3)$ with orbits of lengths 1, 16, and 18 respectively.
- These orbits have been denoted $\{\mathcal{L}\}$, Ψ and Φ
- We consider first the structures obtained from the union of the orbit of length 1 with that of length 18, namely $\{\mathcal{L}\} \cup \Phi$, followed by structures constructed from the orbit of length 16, i.e, Ψ .

Graphs, Designs and Codes from the reprn of degree 35

- Observe that by taking the image of the set $\{\mathcal{L}\} \cup \Phi$, under A_8 we form the blocks of a 1 - $(35, 19, 19)$ design which we denote \mathcal{D}_{19} .
- Since A_8 acts as a rank-3, it follows from Result 1 that the image of Ψ under A_8 defines a strongly regular graph with parameters $(35, 16, 6, 8)$. Denote this graph Γ .
- Equivalently, one could consider the 1 - $(35, 16, 16)$ design, which we denote \mathcal{D}_{16} obtained by orbiting the image of Ψ under A_8 .

Lemma

$\text{Aut}(\mathcal{D}_{19})$, $\text{Aut}(\mathcal{D}_{16})$, and $\text{Aut}(\Gamma)$ are isomorphic to S_8 .

The binary code of Γ

Lemma

- (i) C_{19} is a $[35, 7, 15]_2$ code. Its dual C_{19}^\perp is an *optimal self-orthogonal singly-even* $[35, 28, 4]_2$ code with 840 words of weight 4, and $\mathbf{1} \in C_{19}$.
- (ii) C_Γ is a $[35, 6, 16]_2$ *self-orthogonal doubly-even* code with 35 words of minimum-weight. Moreover $C_\Gamma \subseteq C_{19}$ is a *projective two-weight code*, and C_{19} is a *decomposable \mathbb{F}_2 -module*.
- (iii) C_Γ^\perp is a $[35, 29, 3]_2$ code with 105 words of weight 3, and C_Γ and C_Γ^\perp are optimal codes.
- (iv) $\text{Aut}(C_{19}) = \text{Aut}(C_\Gamma) \cong S_8$.
- (v) S_8 *acts irreducibly on C_Γ as an \mathbb{F}_2 -module*.

Geometry in the codes

- The statements on the parameters of the codes are easily verified.
- Since \mathcal{D}_{19} is the complement of \mathcal{D}_{16} , the difference of any two codewords in C_{16} is in C_{19} .
- As these differences span a subcode of dimension 6 in C_{19} , this subcode must be C_{16} .
- The weight enumerator of C_{19} is as follows

$$W_{C_{19}}(x) = 1 + 28x^{15} + 35x^{16} + 35x^{19} + 28x^{20} + x^{35},$$

and that of C_{16} is given below, denoted $W_{C_{\Gamma}}(x)$.

- Notice from the weight distribution that C_{Γ} is the subcode of C_{19} span by words of weight divisible by four.

Geometry in the codes

- Since \mathcal{D}_{19} is the complement of \mathcal{D}_{16} , the inclusion follows as C_{19} is C_{16} adjoined by the $\mathbf{1}$ vector. So $C_{19} = \langle C_{16}, \mathbf{1} \rangle = C_{16} \oplus \langle \mathbf{1} \rangle$
- Since Γ is a graph that appears in a partition of the symplectic graph $\mathcal{S}_6(2)$, it follows from Peeters [9, Theorem 5.3] that Γ possesses the triangle property and as such it is uniquely determined by its parameters and by the minimality of its 2-rank, which is 6. Thus the dimension of C_Γ is 6.
- The minimum-weight 16 of C_Γ can be deduced using results from Haemers, Peeters and Van Rijkevorseel [7, Section 4.4]. We note that all codewords of C_Γ are linear combinations of at most two rows of the adjacency matrix of Γ .

Geometry in the codes

- Since there are exactly 35 codewords of minimum weight in C_Γ and these correspond to the rows of the adjacency matrix of Γ , these span the code. Now the spanning vectors, have weight 16, so C_Γ is doubly-even and hence self-orthogonal.
- In addition C_Γ is a **two-weight code**, with weight distribution

$$W_{C_\Gamma}(x) = 1 + 35 x^{16} + 28 x^{20}.$$

Since C_Γ^\perp has minimum weight 3 it follows from [Calderbank and Kantor \[2\]](#) that C_Γ is a **projective code**.

- Optimality of C_Γ and C_Γ^\perp follows by [Magma \[1\]](#) and also from the online tables of [Grassl \[6\]](#).
- Note that the 2-modular character table of S_8 is completely known ([Atlas of Brauer Characters](#)) (see [8, 11]) and follows from it that the irreducible 6-dimensional \mathbb{F}_2 -representation is unique.

Strongly regular graphs from the codewords of Γ

- A **two-weight code** is a code which has only two non-zero weights w_1 and w_2 .
- Let w_1 and w_2 (where $w_1 < w_2$) be the weights of a q -ary two-weight code C of length n and dimension k .
- To C we associate a graph $\Lambda(C)$ on q^k vertices as follows: the **vertices** of the graph are identified with the **codewords** and two vertices corresponding to the codewords x and y **are adjacent if and only if** $d(x, y) = w_1$.
- Then $\Lambda(C)$ is a strongly regular graph with parameters (v, k, λ, μ) .
- Following the above, from C_Γ we obtain a strongly regular graph which we denote $\Lambda(C_\Gamma)$ with parameters **$(64, 35, 18, 20)$** and its complement, a strongly regular **$(64, 28, 12, 12)$** graph $\overline{\Lambda(C_\Gamma)}$.

Geometric interpretations

- The words of weight 16 have a geometrical significance: they are the rows of the adjacency matrix of Γ or equivalently the incidence vectors of the blocks of \mathcal{D}_{16} .
- It follows from [Atlas](#) [3] that the objects permuted by the automorphism group are the duads and bisections.
- Moreover, from [Atlas](#) [3] it can also be deduced that the words of weight 16 represent the duads, while those of weight 20, represent the bisections. The stabilizer of a duad is a group isomorphic to $(S_4 \times S_4):2$ while that of a bisection is a group isomorphic to $S_6 \times 2$. Note that these are all maximal subgroups of A_8 and thus A_8 acts primitively on the set of duads and on the set of bisections.

Geometric interpretations

- Viewing A_8 as $L_4(2)$ (the isomorphism could be found in [Dickson](#) and [Taylor](#) [5, 10]) it follows from [Atlas](#) [3] that the objects permuted by the automorphism group are copies of $S_4(2)$ and lines. The codewords of weight 16 represent copies of $S_4(2)$ thereby explaining the connection found in the proof with the symplectic graph $S_6(2)$.
- The codewords of weight 20 represent lines of $L_4(2)$ in this way we can observe the connection established in [Dempwolff](#) [4]. The stabilizer of a copy of $S_4(2)$ is a group isomorphic to $(S_4 \times S_4):2$, while that of a line is a group isomorphic to $S_6 \times 2$. Note that these are all maximal subgroups of A_8 and thus A_8 acts primitively on the set of conjugates of $S_4(2)$ and on the lines.

Geometric interpretations

- The dimension 6 of C_F provides a nice illustration of the isomorphism between A_8 and $\Omega^+(6, 2)$. Therefore using $A_8 \cong \Omega^+(6, 2)$ we can regard the non-zero codewords of C_F as both the non-isotropic and the isotropic points. This in turn indicates that the objects being permuted are the non-isotropic and the isotropic points respectively.
- Finally, the stabilizer of a non-isotropic point under the action of the automorphism group is a maximal subgroup isomorphic with $S_6 \times 2$ while that of an isotropic point is again a maximal subgroup isomorphic to $(S_4 \times S_4):2$.

The ternary code of a 2-(35, 18, 9) design $\bar{\Gamma}$

- We now look at the orbit of length 18, namely Φ . As before, since A_8 acts as a rank-3, it follows from Result 2.1 that the image of Φ under A_8 defines a strongly regular graph with parameters (35, 18, 9, 9). We denote this graph by $\bar{\Gamma}$ where the symbol $\bar{}$ is standard for denoting the complement of Γ .
- Notice that $\bar{\Gamma}$ is 2-(35, 18, 9) design
- Since the order of $\bar{\Gamma}$ is 9 the only codes of interest are ternary.
- We examine the codes obtained from the ternary row span of the adjacency matrix of $\bar{\Gamma}$.

Lemma

- (i) $C_{\bar{\Gamma}}$ is a $[35, 13, 12]_3$ code, $C_{\bar{\Gamma}}^{\perp}$ is a $[35, 22, 5]_3$ with 112 words of weight 5, and $\mathbf{1} \in C_{\bar{\Gamma}}^{\perp}$
- (ii) $\text{Aut}(\bar{\Gamma}) = \text{Aut}(C_{\bar{\Gamma}}) \cong S_8$.

A self-dual $[72, 36, 8]_2$ code from $\bar{\Gamma}$

- Let A be the incidence matrix of $\bar{\Gamma}$, and $A^+ = \begin{pmatrix} A & \mathbf{1}^t \\ \mathbf{1} & 0 \end{pmatrix}$ where $\mathbf{1}$ is the all one vector of length 35.
- A generator matrix of a double-even self-dual code of length 72 can be obtained as $\begin{pmatrix} A^+ & I_{36} \end{pmatrix}$.
We used this method to construct a $[72, 36, 8]_2$ formally self-dual code denoted \mathcal{T} , from the incidence matrix of $\bar{\Gamma}$.

A self-dual $[72, 36, 8]_2$ code from $\bar{\Gamma}$

Corollary

The binary code \mathcal{T} of $(A^+ I_{36})$ is a self-dual doubly even $[72, 36, 8]_2$ code, with automorphism group isomorphic to $2^{15}:S_6(2)$.

- The weight enumerator of \mathcal{T} is as follows:

$$\begin{aligned}W_{\mathcal{T}}(x) = & 1 + 945 x^8 + 30576 x^{12} + 535932 x^{16} + 17267040 x^{20} \\ & + 455965020 x^{24} + 4438423440 x^{28} + 16506508662 x^{32} \\ & + 25882013504 x^{36} + 16506508662 x^{40} \\ & + 4438423440 x^{44} + 455965020 x^{48} + 17267040 x^{52} \\ & + 535932 x^{56} + 30576 x^{60} + 945 x^{64} + x^{72}.\end{aligned}$$

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