

Network Codes and Designs over Finite Fields

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- Designs
- Network Codes
- Construction

Designs



- a set of v points

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- a set of blocks (block = set of points)

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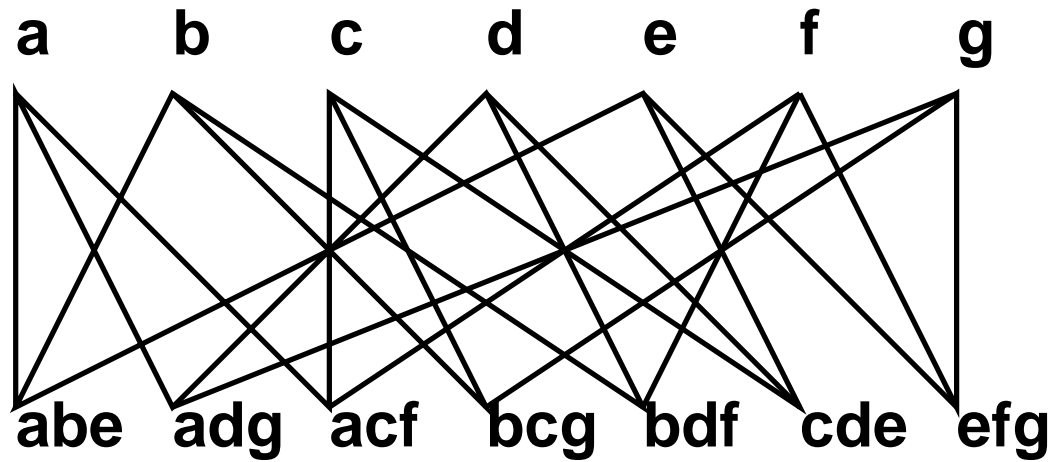
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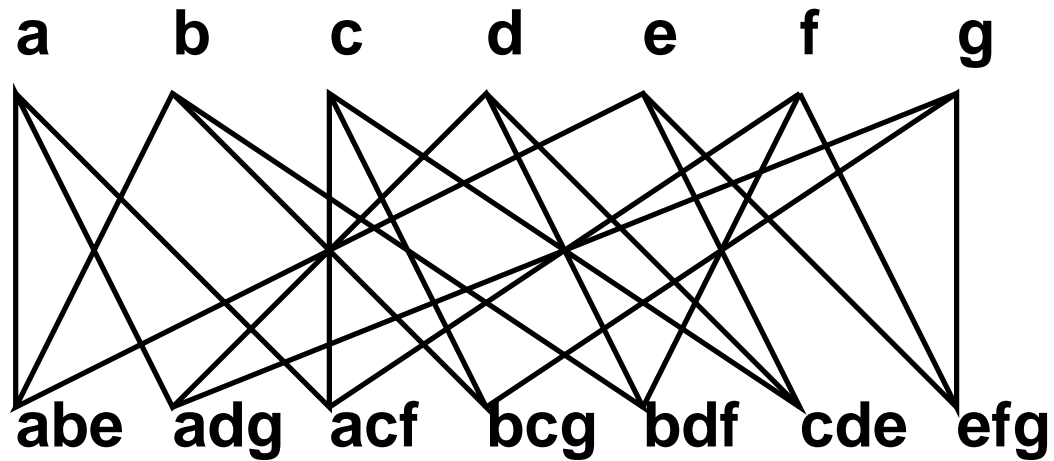
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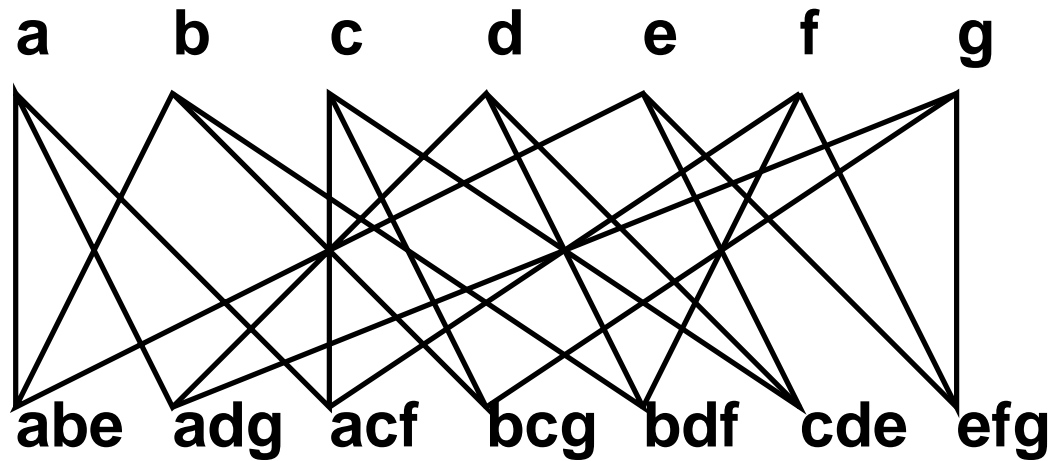
2 - (7, 3, 1) design



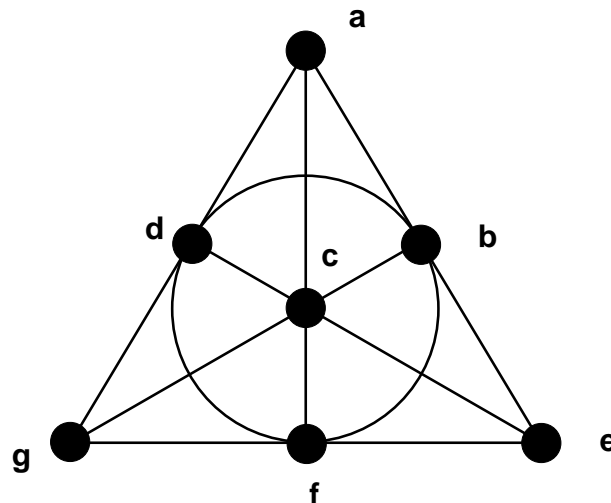
Heawood graph



Heawood graph



Fano plane



Designs over Finite Fields

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linear v -space $GF(q)^v$

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- $t - (v, k, \lambda) q$ -Design

each t -space of $GF(q)^v$ is in exactly λ of the k -spaces

known:

- Thomas (1987): first to study, 2–designs
- Braun, Kerber, Laue (2005): first 3–design

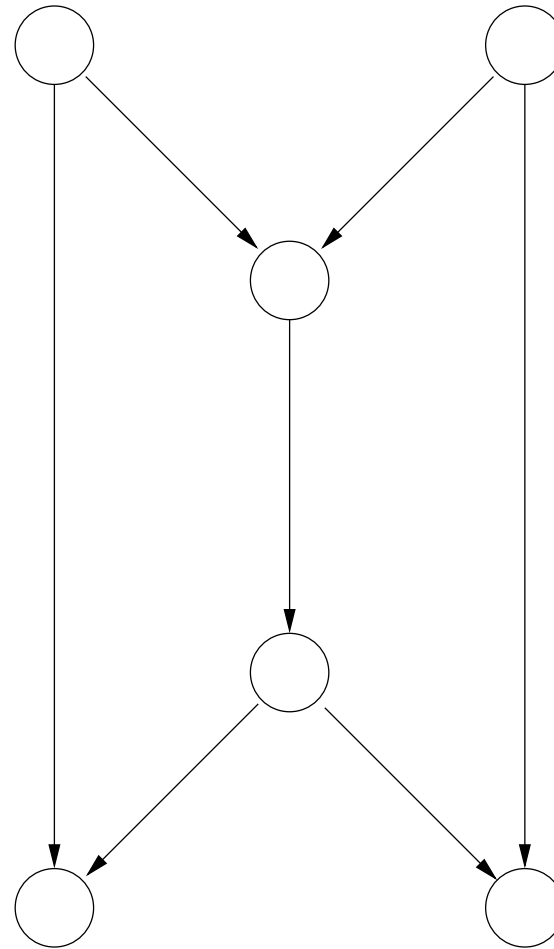
open problems:

- q –analog of the Fano plane?
- Steiner systems ? ($\lambda = 1$)
- $t > 3$?

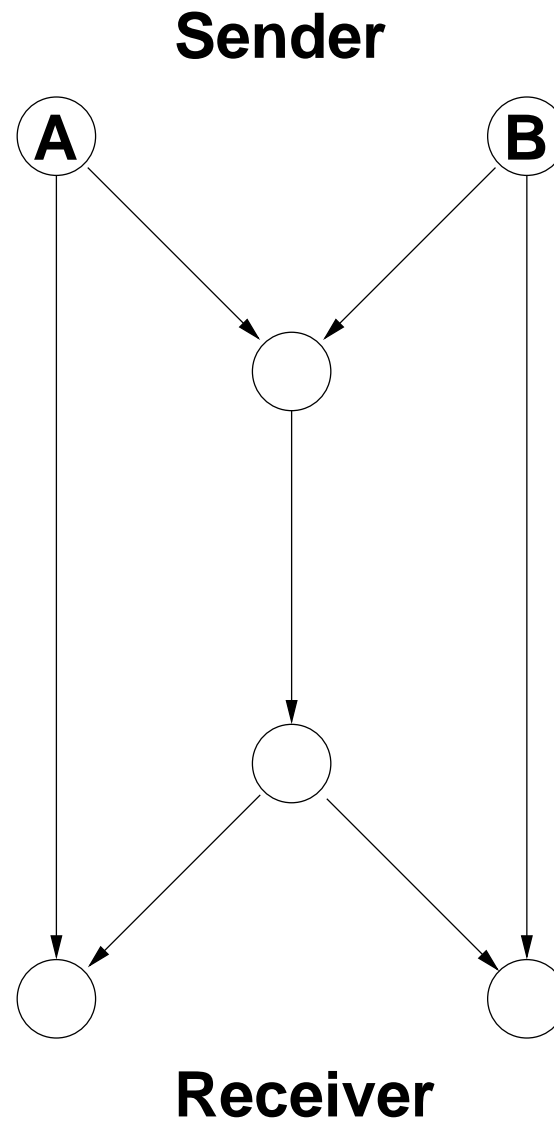
Network Codes

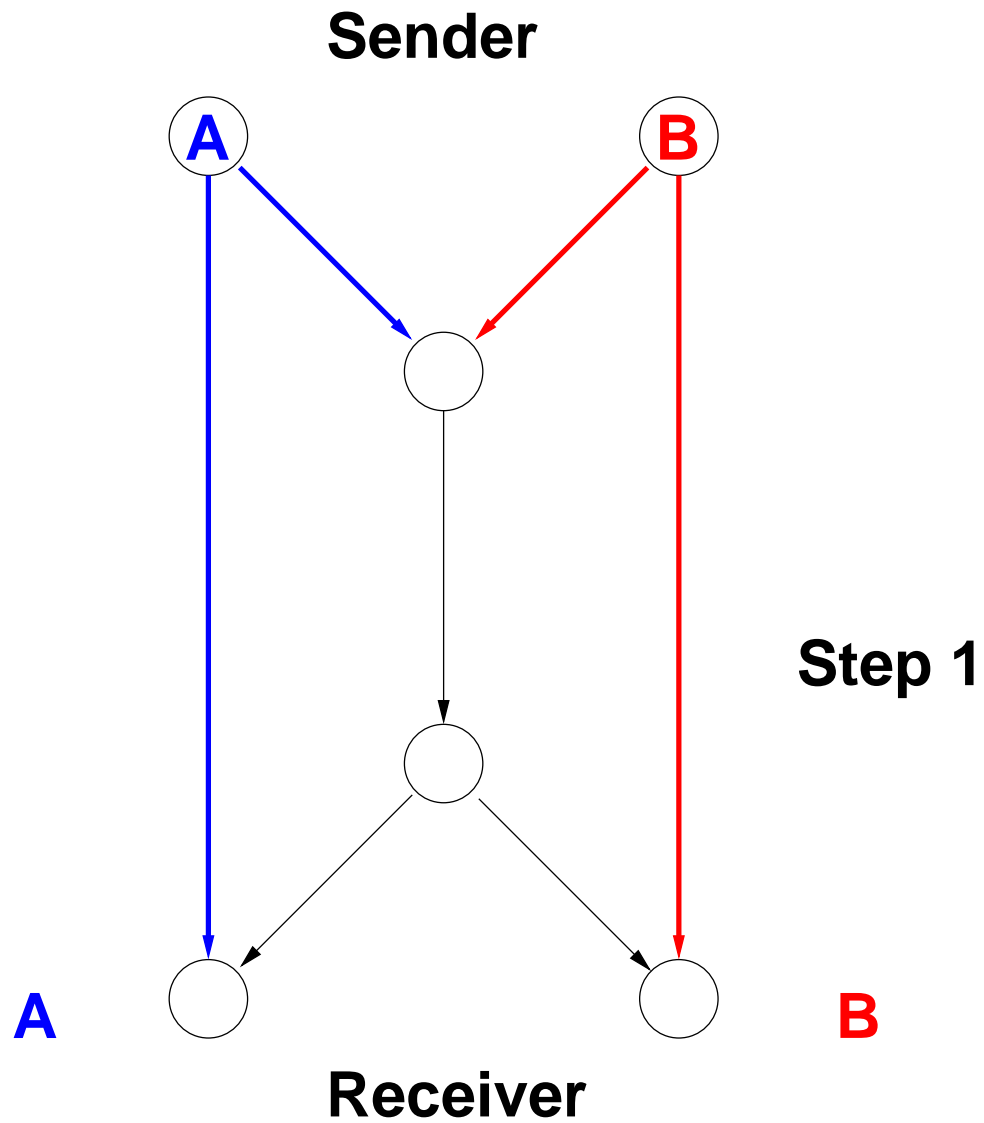


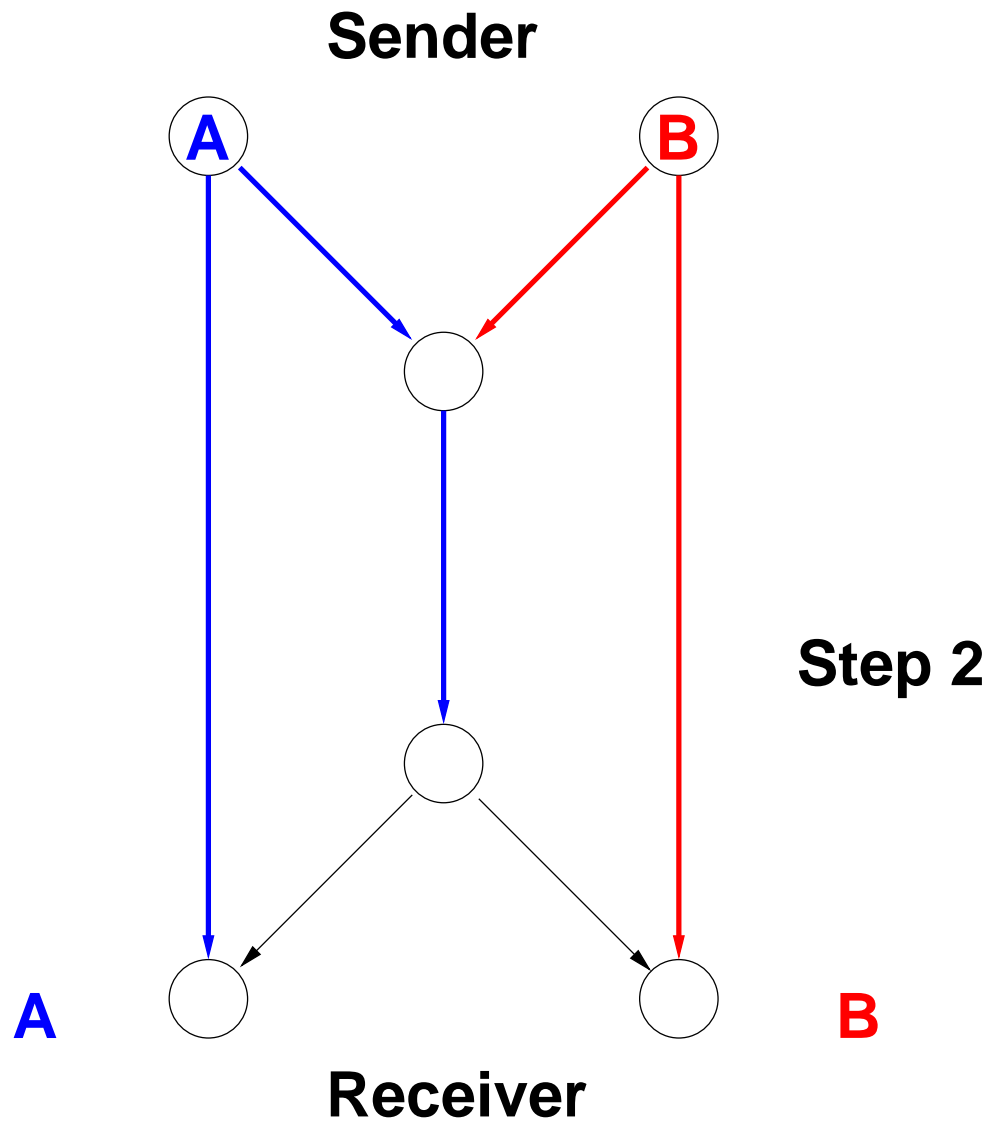
Sender

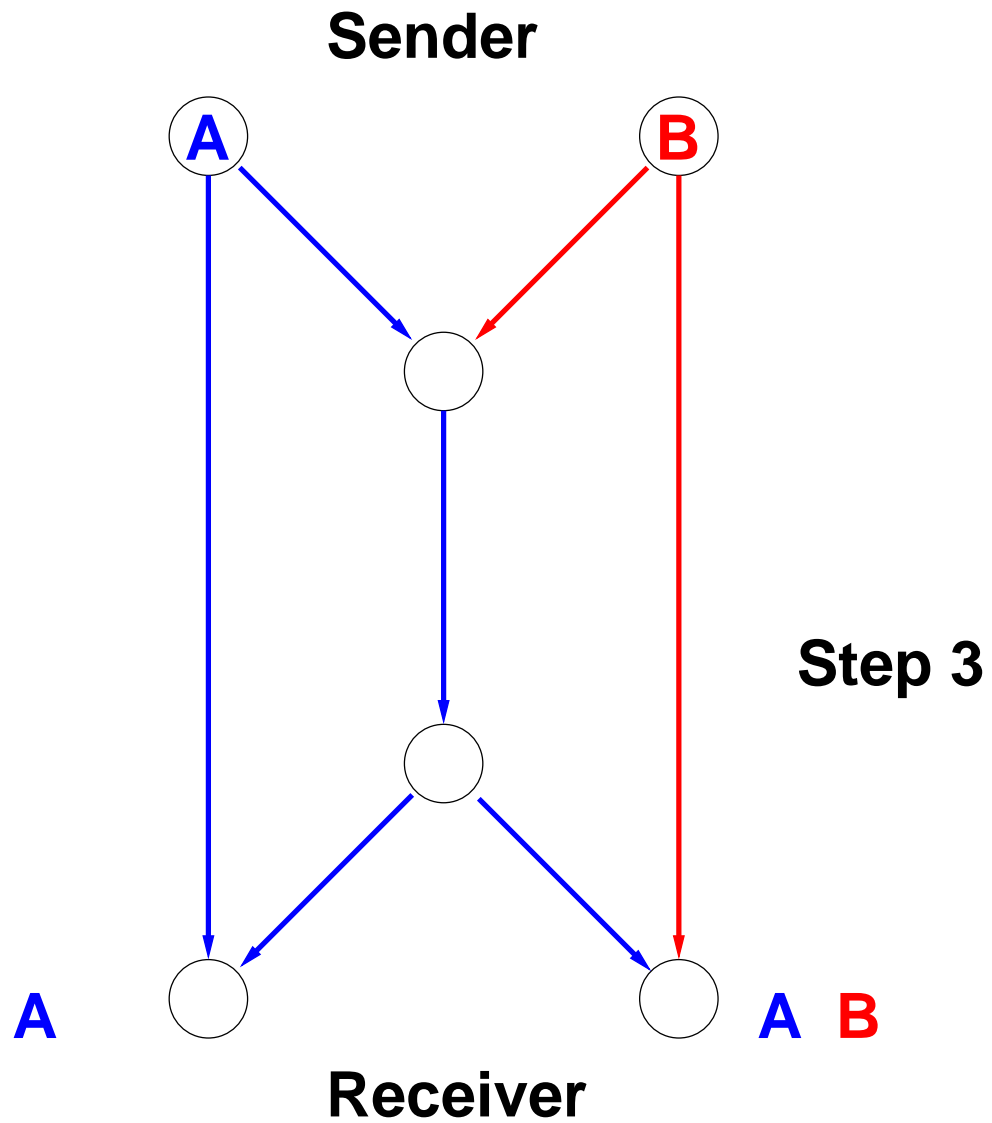


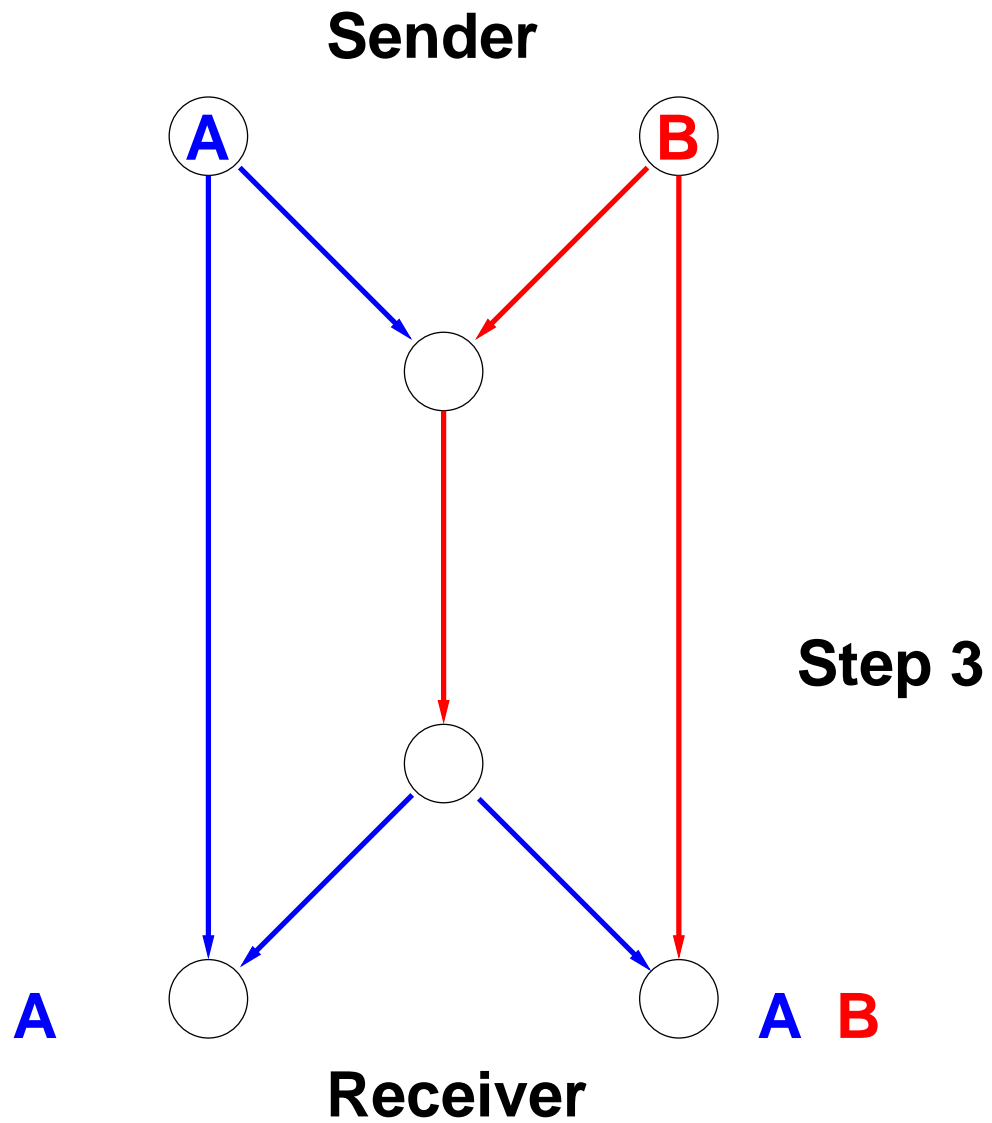
Receiver

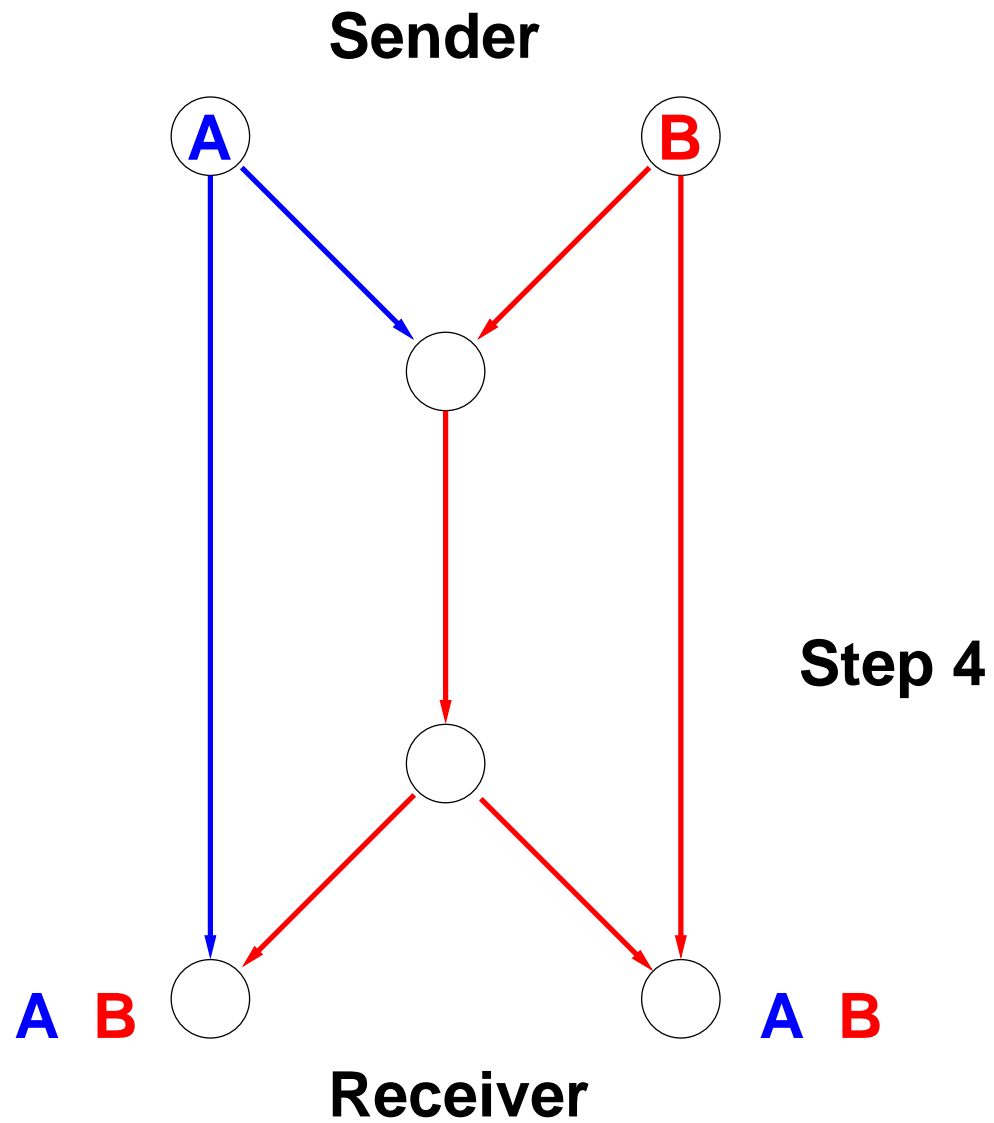


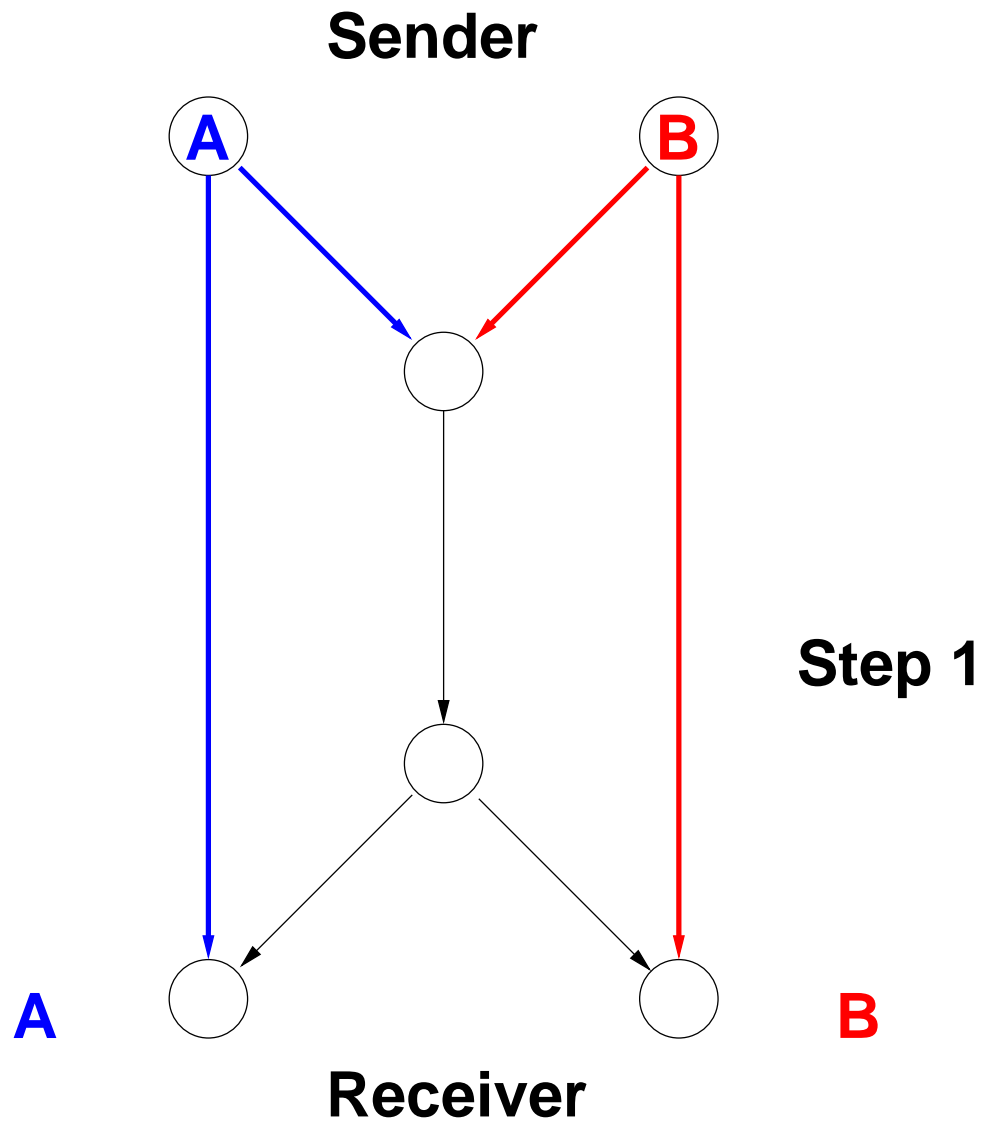


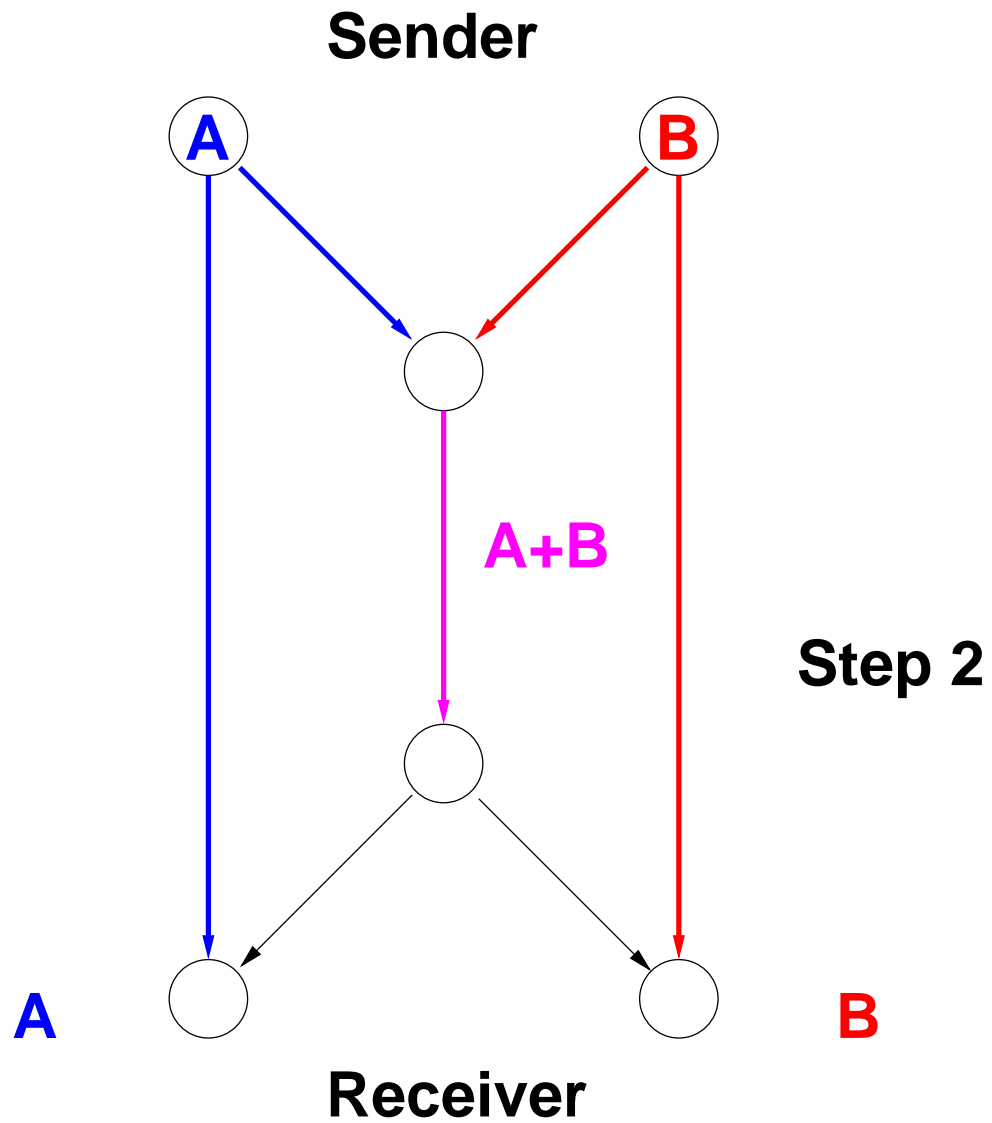


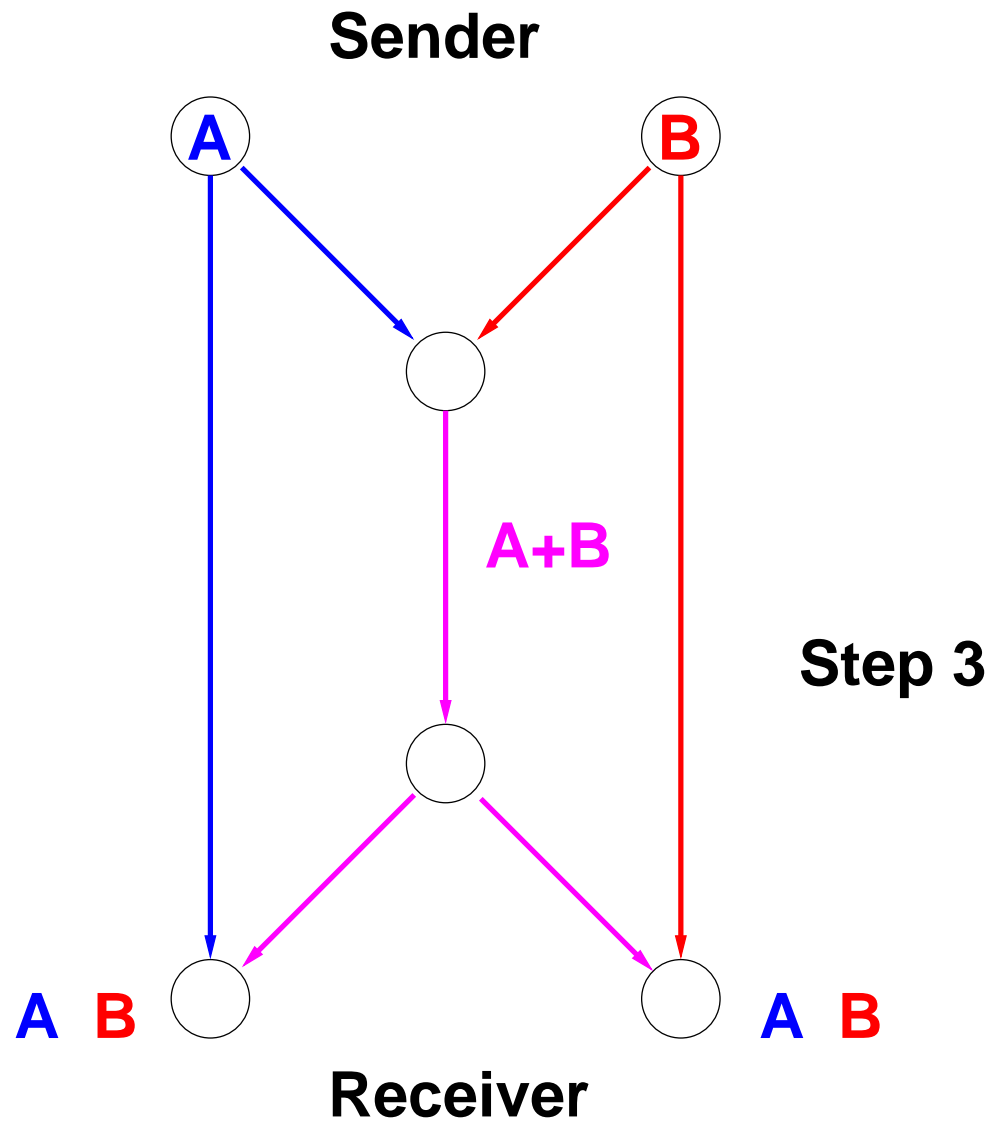












message:

- linear space

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single node:

- receives vectors
- sends some linear combination of the incoming vectors

Error-Correcting Network Codes

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- linear subspace of $GF(q)^v$

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U, W subspace of $GF(q)^v$:

$$d(U, W) = \dim(U) + \dim(W) - 2\dim(U \cap W)$$

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constant dimension codes \approx constant weight codes

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Define $A_q(v, k, d)$ as the maximal size (= number of codewords) of a constant dimension code with minimum distance d , dimension of codewords = k , and ambient space = $GF(q)^v$

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open problems:

- find lower and upper bounds for $A_q(v, k, d)$
- find constructions of 'good' codes
- special case $A_2(7, 3, 4) = \text{Fano plane}$

Construction



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$D :=$ incidence matrix between k -spaces and t -spaces in $GF(q)^v$

$$D_{U,V} := \begin{cases} 1 & t\text{-space } U \text{ is subspace of } k\text{-space } W \\ 0 & \text{else} \end{cases}$$

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- $D^G :=$ shrunked matrix

\Rightarrow number of columns = number of orbits on k -spaces
number of rows = number of orbits on t -spaces

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solution = network code with prescribed automorphisms and minimum distance $2(k - t + 1)$.

Results (binary)

v	k	number of codewords: new	old	d
6	3	77	71	4
7	3	304	294	4
8	3	1275	1164	4
9	3	5621	4657	4
10	3	21483	18631	4
11	3	79833	74531	4
12	3	315315	298139	4



A. Kohnert, S. Kurz: *Construction of Large Constant Dimension Codes With a Prescribed Minimum Distance*, LNCS, 2008.

R. Kötter, F. Kschischang: *Coding for errors and erasures in random network coding*, IEEE Transactions on Information Theory, **54**, 3579–3590, 2008.

Thank you very much for your attention.