Codes and Sequences Over Finite Rings

Eimear Byrne

Codes and Sequences Over Finite Rings

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ALCOMA10

Outline

Codes and Sequences Over Finite Rings

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- Background
- Rings and Weights
- Sequences and Codes
- Examples

A Binary Code

Codes and Sequences Over Finite Rings

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Let
$$(n, s) = 1$$
 and let $d = 2^s + 1$. Consider the binary code: $C = \{c_{\alpha,\beta}(x) = \text{Tr}(\alpha x) + \text{Tr}(\beta x^d), \alpha, \beta \in GF(2^n)\}.$

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C has generator matrix

$$\left[\begin{array}{c|c} x_1 & x_2 & \cdots & x_{2^n-1} \\ x_1^d & x_2^d & \cdots & x_{2^n-1}^d \end{array}\right],$$

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and

$$w_H(c_{\alpha,\beta}) = \left(2^n - \sum_{x \in GF(2^n)} (-1)^{\operatorname{Tr}(\alpha x) + \operatorname{Tr}(\beta x^d)}\right)/2.$$

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C has length $2^n - 1$. For odd n it has dimension 2n and 3 non-zero weights:

$${2^{n-1}-2^{\frac{n-1}{2}},\ 2^{n-1},\ 2^{n-1}+2^{\frac{n-1}{2}}}.$$

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For a finite ring R, $\hat{R}:=\operatorname{Hom}_{\mathbb{Z}}(R,\mathbb{C}^{\times})$, is an R-R bimodule:

$$^{r}\chi(x) = \chi(rx), \quad \chi^{r}(x) = \chi(xr)$$

for all $x, r \in R, \chi \in \hat{R}$.

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The following are equivalent definitions:

- \blacksquare R is a Frobenius ring
- soc_RR is left principal,
- \blacksquare $_R(R/rad\ R) \simeq soc_R R$,
- $RR \simeq R\hat{R}$

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Then $_R\hat{R}=_R\langle\chi\rangle$ for some (left) generating character χ .

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The following are examples of Frobenius rings.

- integer residue rings \mathbb{Z}_m
- any semi-simple ring
- principal ideal rings
- direct products of Frobenius rings
- matrix rings over Frobenius rings
- group rings over Frobenius rings

Homogeneous Weights

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Definition

A weight $w: R \longrightarrow \mathbb{Q}$ is (left) homogeneous, if w(0) = 0 and

- 1 If Rx = Ry then w(x) = w(y) for all $x, y \in R$.
- **2** There exists a real number γ such that

$$\sum_{y \in Rx} w(y) \ = \ \gamma \, |Rx| \qquad \text{for all } x \in R \setminus \{0\}.$$

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Example

On every finite field \mathbb{F}_q the Hamming weight is a homogeneous weight of average value $\gamma = \frac{q-1}{q}$.

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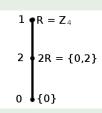
Example

On every finite field \mathbb{F}_q the Hamming weight is a homogeneous weight of average value $\gamma = \frac{q-1}{q}$.

Example

On \mathbb{Z}_4 the Lee weight is homogeneous with $\gamma = 1$.

X	0	1	2	3
$w_{\text{Lee}}(x)$	0	1	2	1



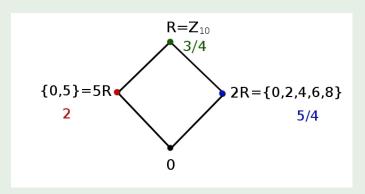
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Example

On \mathbb{Z}_{10} the following weight is homogeneous with $\gamma = 1$:

X	0	1	2	3	4	5	6	7	8	9
$w_{hom}(x)$	0	<u>3</u>	<u>5</u>	<u>3</u>	<u>5</u>	2	<u>5</u>	<u>3</u>	<u>5</u>	<u>3</u>



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Example

On a local Frobenius ring R with q-element residue field the weight

$$w:R\longrightarrow \mathbb{R},\quad x\mapsto \left\{ egin{array}{ll} 0 & : & x=0, \\ rac{q}{q-1} & : & x\in soc(R), \ x
eq 0, \\ 1 & : & \text{otherwise}, \end{array}
ight.$$

is a homogeneous weight of average value $\gamma = 1$.

Homogeneous Weights of FFRs

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Theorem (Honold)

Let R be a finite Frobenius ring with generating character χ . Then the homogeneous weights on R are precisely the functions

$$w: R \longrightarrow \mathbb{R}, \quad x \mapsto \gamma \Big[1 - \frac{1}{|R^{\times}|} \sum_{u \in R^{\times}} \chi(xu) \Big]$$

where γ is a real number.

Characters and Trace Maps

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Let R > S be Frobenius rings.

Definition

Let T be an S-module epimorphism $T: {}_SR \longrightarrow {}_SS$ whose kernel contains no non-trivial left ideal of R. We say that T is a trace map from R onto S.

Characters and Trace Maps

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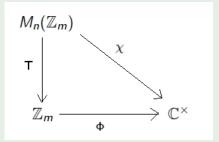
A generating character $\Phi \in \hat{S}$ determines a generating character $\chi \in \hat{R}$ as:

$$\chi(x) = \Phi(T(x)) \ \forall \ x \in R.$$

Codes and Sequences Over Finite Rings

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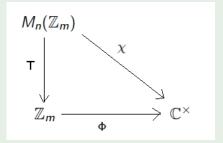
Example (Characters and Traces on $M_n(\mathbb{Z}_m)$)



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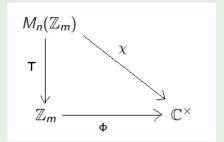


 $\Phi(x) = \omega^x, \omega$ a primitive *m*th root of unity in \mathbb{C}^{\times}

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Example (Characters and Traces on $M_n(\mathbb{Z}_m)$)



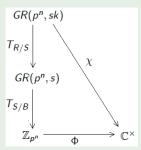
- $\Phi(x) = \omega^x, \omega$ a primitive *m*th root of unity in \mathbb{C}^{\times}
- T is the usual trace map from $M_n(\mathbb{Z}_m)$ onto \mathbb{Z}_m .

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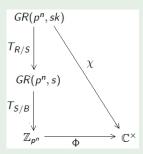
Example (Characters and Traces on Galois Rings)

Let $R = GR(p^n, sk), S := GR(p^n, s), B := \mathbb{Z}_{p^n}$.



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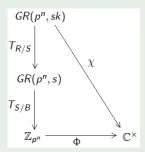
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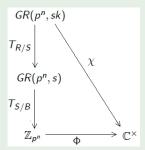


- $\Phi(x) = \omega^x, \omega$ a primitive p^n th root of unity in \mathbb{C}^{\times}
- $\sigma: R \longrightarrow R: \sum_{i=0}^{n} p^{i} a_{i} \mapsto \sum_{i=0}^{n} p^{i} a_{i}^{p} \in Aut(R)$

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Example (Characters and Traces on Galois Rings)

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- $\sigma: R \longrightarrow R: \sum_{i=0}^{n} p^{i} a_{i} \mapsto \sum_{i=0}^{n} p^{i} a_{i}^{p} \in Aut(R)$
- \blacksquare $T_{R/S}: R \longrightarrow S: a \mapsto a + \sigma^{s}(a) + \cdots + \sigma^{s(k-1)}(a)$

A Subring Subcode

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For any map $f: R \longrightarrow R$, we define the left *S*-linear subring subcode

$$C_f = \{c_{\alpha,\beta}^f : R \longrightarrow S : x \mapsto T(\alpha x + \beta f(x)) : \alpha, \beta \in R\}.$$

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$$\begin{split} w(c_{\alpha,\beta}^f) &= \sum_{x \in R} w(c_{\alpha,\beta}^f(x)) \\ &= |R| - \frac{1}{|S^{\times}|} \sum_{u \in S^{\times}} \sum_{x \in R} \Phi^u(T(\alpha x + \beta f(x))) \end{split}$$

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Definition

Let R > S be Frobenius rings with trace map $T : {}_{S}R \longrightarrow {}_{S}S$. Let $f : R \longrightarrow R$. For each $\alpha, \beta \in R$, define

$$W^f(\alpha,\beta) := \frac{1}{|S^{\times}|} \sum_{u \in S^{\times}} \sum_{x \in R} \chi^u(\alpha x + \beta f(x)) = |R| - w(c_{\alpha,\beta}^f).$$

The spectrum of f is the set

$$\Lambda_f := \{ W^f(\alpha, \beta) : \alpha, \beta \in R \}.$$

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- If $|\Lambda_f| = k + 1$ then C_f has exactly k non-zero weights.
- One of the weights of C_f is |R|.

Frank Sequences

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Theorem

Let $R = S = GR(p^2, r)$, p prime. Write $a = a_0 + pa_1$ for each $a \in R$. Let

$$f:R\longrightarrow R:a\mapsto pa_0a_1.$$

Then

$$\Lambda_f = \{p^{2r}, p^r, 0\}.$$

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• C_f has length $p^{2r} - 1$, size p^{3r} and weight enumerator

$$1+p^{r}(p^{r}-1)X^{p^{r}(p^{r}-1)}+(p^{r}-1)(p^{2r}+1)X^{p^{2r}}.$$

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- If we let $S = \mathbb{Z}_{p^n}, r > 1$ then

$$\Lambda_f = \{p^{2r}, p^r, -\frac{p^r}{p-1}, 0\}.$$

Chu Sequences

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Theorem

Let
$$R = S = \mathbb{Z}_{2p}$$
, p prime. Let

$$f:R\longrightarrow R:a\mapsto a^2.$$

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Theorem

Let $R = S = \mathbb{Z}_{2p}$, p prime. Let

$$f:R\longrightarrow R:a\mapsto a^2.$$

Then

$$\Lambda_f = \{2p, \frac{2p}{p-1}, 0\}.$$

 C_f has length 2p-1, size $2p^2$ and weight enumerator

$$1 + (1 + 4(p-1) + (p-1)^2)X^{2p} + (p-1)^2X^{2p\frac{p-2}{p-1}}$$
.

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- Then each element $a \in R$ can be expressed as

$$a = a_m + a_t$$

for some unique $a_m \in M$, $a_t \in T$, where $T \setminus \{0\}$ is a cyclic subgroup of order $|K^{\times}|$ in R^{\times} .

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This decomposition can be useful for evaluating the spectrum of a function.

Compatibility of χ with Aut(R)

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Suppose that

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Then

$$\chi(f(a)) = \chi(\sigma(a)a - \sigma(a_m)a_m)$$

$$= \chi(\sigma(a_t)a_t - \sigma(a_m)a_t - \sigma(a_t)a_m)$$

$$= \chi(\sigma(a_t)a_t)\chi((\sigma^{-1}(a_t) - \sigma(a_t))a_m).$$

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Theorem

Let R be a finite local commutative Frobenius ring. Let $\sigma \in Aut(R)$ satisfy $\chi(\sigma(x)) = \chi(x)$ for all $x \in R$. Define

$$f:R\longrightarrow R:a\mapsto \sigma(a)a-\sigma(a_m)a_m.$$

Then

$$\Lambda_f = \{ |R|, |M|, \frac{|R||M|}{|R^{\times}|}, 0 \}.$$

Theorem

Let R be a finite local commutative Frobenius ring. Let $\sigma \in Aut(R)$ satisfy $\chi(\sigma(x)) = \chi(x)$ for all $x \in R$. Define

$$f: R \longrightarrow R: a \mapsto \sigma(a)a - \sigma(a_m)a_m$$
.

Then C_f has length |R|-1 and non-zero weights

$$\{|R|,|R|-|M|,|R|(1-\frac{|M|}{|R^{\times}|})\}.$$

More

Codes and Sequences Over Finite Rings

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- Find more functions on local rings that give codes with small spectra.
- Determine functions that yield 2-weight codes (especially modular or projective regular codes).
- Nonlinearity.

Applications - Strongly Regular Graphs

Codes and Sequences Over Finite Rings

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